

lecture 16. Orthogonality and Least Squares Problems

$u, v \in \mathbb{R}^m$, $u \perp v$ if $u^T v = 0$.

In this case, $\|u+v\|^2 = \|u\|^2 + \|v\|^2$.

Two subspaces S and T : $S \perp T$ if

$\forall s \in S$, and $\forall t \in T$, $s^T t = 0$, i.e., $S \perp t$.

Four subspaces associated with $A \in \mathbb{R}^{m \times n}$

1. If $S \perp T$, then $S \cap T = \{0\}$

The formula $\dim(S \cap T) + \dim(S \perp T) = \dim(S) + \dim(T)$

implies $\dim(S) + \dim(T) = \dim(S \cap T) \leq m$

So two planes in \mathbb{R}^3 can't be orth. to each other

2. When $S \perp T$, and $\dim(S) + \dim(T) = m$.

then T is called the orthogonal complement of S

denoted by $S^\perp = T$, Vers versa, $T^\perp = S$.

and $(S^\perp)^\perp = S$.

$A \in \mathbb{R}^{m \times n}$

3. $C(A)^\perp = N(AT)$, $\forall x \in \mathbb{R}^m$. $\left\{ \begin{array}{l} \text{Example} \\ ATAx = 0 \Rightarrow Ax = 0 \\ Ax \in N(AT) \\ C(A) \Rightarrow Ax = 0 \end{array} \right.$

$x = x_c + x_n$, with $x_c \in C(A)$, and $x_n \in N(AT)$

and $x_c \perp x_n$.

$$Ax = 0 \Leftrightarrow \begin{pmatrix} A^T x \\ \vdots \\ A_{m \times n}^T x \end{pmatrix} = 0$$

$C(A^T)^{\perp} = N(A)$, $\forall y \in \mathbb{R}^n$

$y = y_c + y_n$, $y_c \in C(AT)$, $y_n \in N(A)$, $y_c \perp y_n$

Example. Assume $Ax = b$ has solution, find the

solution with minimal length $\| \cdot \|$. Any solution

$x = a + b$, with $a \in C(AT)$, $b \in N(A)$

$$\|x\|^2 = \|a\|^2 + \|b\|^2, \text{ set } b = 0.$$

Example $S \subset \mathbb{R}^3$, and $S = \{v_0\}$, $S^\perp = ?$

If $S = \text{Span}\left\{\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}\right\}$, $S^\perp = \text{Span}\left\{\begin{pmatrix} 1 \\ -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ -1 \end{pmatrix}\right\}$

$S = \text{Span}\left\{\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ -1 \\ 1 \end{pmatrix}\right\}$, $S^\perp = \text{Span}\left\{\begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}\right\}$

Example. $a \perp \text{Span}\{v_1, \dots, v_k\}$ iff. $a \perp v_i, i=1, \dots, k$

Proof. " \Rightarrow " Since $v_i \in \text{Span}\{v_1, \dots, v_k\} \Rightarrow a \perp v_i$.

" \Leftarrow " If $a \in \text{Span}\{v_1, \dots, v_k\}$, $a = \sum_{i=1}^k \alpha_i v_i$

$$a^T a = \sum_{i=1}^k \alpha_i a^T v_i = 0$$

least squares problems.

$$A \in \mathbb{R}^{m \times n}$$

We have consider $Ax = b$, when $b \in C(A)$, there might be one or many solutions. But if $b \notin C(A)$, we can ask for an approximate solution in the sense of making $r = b - Ax$ small but not exactly zero, r is the so-called residual of an approximate sol x .

least squares is one approach of quantifying the size of "smallness" of r , in terms of its

$$\text{L}^2\text{-norm} \quad \|r\|^2 = r_1^2 + r_2^2 + \dots + r_m^2, \quad r = \begin{pmatrix} r_1 \\ \vdots \\ r_m \end{pmatrix}$$

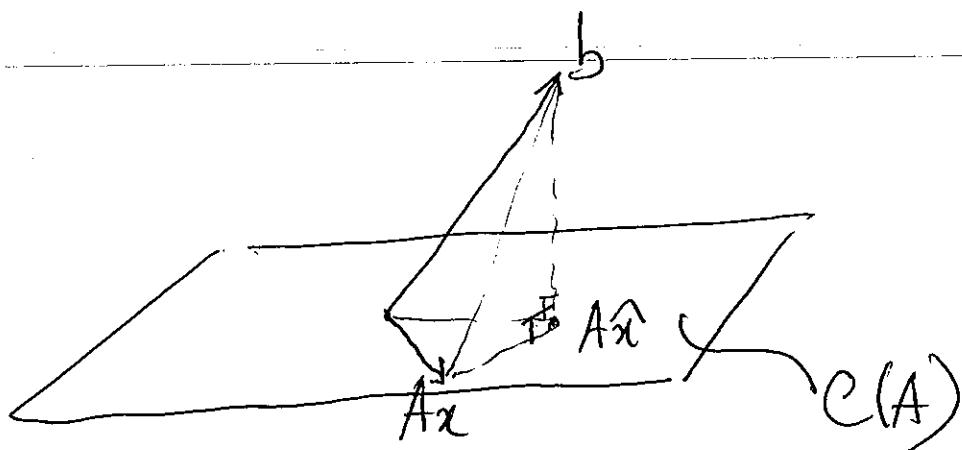
Other possibilities exist, e.g. $r_i = (b_i - A_i x)^2$

$$\|r\|_1 = |r_1| + |r_2| + \dots + |r_m| \quad A = \begin{pmatrix} A_1 \\ \vdots \\ A_m \end{pmatrix}$$

So we want find an $x \in \mathbb{R}^n$, s.t.

$$\min_{x \in \mathbb{R}^n} \|r\|^2 = \min_{x \in \mathbb{R}^n} \|b - Ax\|^2$$

This is to say we find $\hat{Ax} \in C(A)$ s.t. the residue $b - \hat{Ax}$ is the smallest in $\|\cdot\|_2$



Orthogonality principle: $\|r\|^2$ is minimized iff.

$$r \perp C(A) \Leftrightarrow a_i \perp r, \forall i=1, \dots, n.$$

where $A = (a_1, \dots, a_n) \in \mathbb{R}^{m \times n}$

From $\begin{pmatrix} a_1^T r \\ a_2^T r \\ \vdots \\ a_n^T r \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$ we obtain

$$A^T r = 0 \quad A^T(b - A\hat{x}) = 0$$

Normal Eq $A^T A \hat{x} = A^T b$, this is an $n \times n$ linear

system. Since $\text{rank}[A^T A, A^T b] = \text{rank}(A^T A)$

it always has a solution, no matter what b is.

It has a unique solution, $\hat{x} = (A^T A)^{-1} A^T b$

when A has full column rank, i.e., $\text{r}(A) = n \leq m$.

$$\text{Now } p = A\hat{x} = A(A^T A)^{-1} A^T b = P_A b$$

P_A is a projection matrix. $P = P^T, P^2 = P$.

Will discuss $\text{r}(A) < n$ later!

$$\text{Now } b = P_A b + (I - P_A) b$$

Since $P_A b \perp (I - P_A) b$. for arbitrary $x \in \mathbb{R}^n$

$$\|b - Ax\|^2 = \|(P_A b - Ax) + (I - P_A) b\|^2$$

Now $Ax \in C(A)$, $Ax \perp (I - P_A) b$.

$$= \|P_A b - Ax\|^2 + \|(I - P_A) b\|^2$$

The residual is minimized iff $Ax = P_A b$.

Given two lines :

$$P = (x, x, x), \quad x \in \mathbb{R}.$$

$$Q = (y, 3y, -1), \quad y \in \mathbb{R}.$$

Do the two lines meet at some point ?

We can formulate the above as a LS problem.

$$\begin{aligned} \min_{P, Q} \|P - Q\|^2 &= \min_{x, y} \left\| \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}x - \begin{pmatrix} 1 \\ 3 \\ 0 \end{pmatrix}y - \begin{pmatrix} 0 \\ 0 \\ -1 \end{pmatrix} \right\|^2 \\ &= \left\| \begin{pmatrix} 1 & 1 \\ 1 & 3 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} - \begin{pmatrix} 0 \\ 0 \\ -1 \end{pmatrix} \right\|^2. \end{aligned}$$

Example. Fitting a straight line

$$f(t) = c + dt$$

to a set of points (t_i, b_i) , $i=1, \dots, m$.

The linear system is

$$Ax = b, \text{ with } A = \begin{pmatrix} 1 & t_1 \\ \vdots & \vdots \\ 1 & t_m \end{pmatrix}, \quad \begin{array}{l} c + dt_1 = b_1 \\ c + dt_2 = b_2 \\ \vdots \\ c + dt_m = b_m \end{array}$$

$$b = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix}$$

$$ATA = \begin{pmatrix} m & \sum t_i \\ \sum t_i & \sum t_i^2 \end{pmatrix}, \quad ATb = \begin{pmatrix} \sum b_i \\ \sum t_i b_i \end{pmatrix}$$

$$\hat{x} = (ATA)^{-1}ATb = \frac{1}{m\sum t_i^2 - (\sum t_i)^2} \begin{pmatrix} \sum t_i^2 - \bar{t}_i \\ -\bar{t}_i \end{pmatrix} \begin{pmatrix} \sum b_i \\ \sum t_i b_i \end{pmatrix}$$

$$= \frac{\left(\begin{pmatrix} \sum t_i^2 \end{pmatrix} \left(\sum b_i \right) - \left(\sum t_i \right) \left(\sum t_i b_i \right) \right)}{\left(m \left(\sum t_i b_i \right) - \left(\sum t_i \right) \left(\sum b_i \right) \right)} / \left(m\sum t_i^2 - (\sum t_i)^2 \right)$$

This can be easily generalized to fitting a polynomial to a set of data, or supervised learning.

Supervised learning.

Let $g_1(x), \dots, g_n(x)$ be a set of basis functions, e.g., $1, x, \dots, x^{n-1}$, $x \in \mathbb{R}$, but in general $x \in \mathbb{R}^p$, multi-dimensional.

We are given a set of data points (x_i, y_i)

We want find a linear combination of the f_i 's

$$g(x) = \alpha_1 g_1(x) + \dots + \alpha_n g_n(x)$$

s.t. $g(x_i) \approx y_i$, $i=1, \dots, m$. If we minimize $\sum_{i=1}^m (y_i - g(x_i))^2$, we have the

LS problem $\|Ax - b\|^2$, where

$$A = \begin{pmatrix} g_1(x_1) & \dots & g_n(x_1) \\ \vdots & & \vdots \\ g_1(x_m) & \dots & g_n(x_m) \end{pmatrix}, \quad x = \begin{pmatrix} \alpha_1 \\ \vdots \\ \alpha_n \end{pmatrix}, \quad b = \begin{pmatrix} y_1 \\ \vdots \\ y_m \end{pmatrix}.$$

Th let $C(A)$ denote the col space of A .

Thm If $C(A)$ is a subspace of $C(B)$, then $r(A) \leq r(B)$.

Let's prove the following $r(A+B) \leq r([A, B])$

Proof $A = [a_1, \dots, a_n]$, $B = [b_1, \dots, b_n]$.

$$a_i \in C([A, B]), b_i \in C([A, B])$$

hence $a_i + b_i \in C([A, B])$, which implies

$$C(A+B) \subseteq C(C([A, B])) \Rightarrow r(A+B) \leq r([A, B])$$

Corollary $r(A+B) \leq r(A) + r(B)$

Corollary $|r(A) - r(B)| \leq r(A-B)$

$A \in \mathbb{R}^{n \times n}$, $B \in \mathbb{R}^{n \times p}$ if $AB=0$, then

$$r(A) + r(B) \leq n.$$

Proof. $AB=0 \Rightarrow C(B) \subseteq N(A) \Rightarrow r(B) \leq \dim(N(A)) = n - r(A)$